ch22: Graph Algorithms

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Outline for today

Depth-first search

- Parenthesis structure
- Edge classification
- Topological sort
- Finding strongly-connected components

Semester overview



Breadth-first search algorithm

Input: vertex list, adjacency list (linked lists), start
Output: modify vertex list to add parent pointers

→ def BFS(V, E, start):

- initialise all vertices to be white, with NULL parent
 - initialise start to be grey
 - initialise FIFO: Q.push(start)
 - while Q.notempty():
 - u = Q.pop()
 - for each v in E.adj[u]:
 - if v.colour == white:
 - v.colour = grey
 - v.parent = u
 - Q.push(v)
 - u.colour = black

• Complexity: O(V + E)

Е

Depth-first search

Explore as deep as we can go Backtrack to explore other paths Recursive algorithm Colouring: white = undiscovered Grey = discovered Black = finished (visited all neighbours) Add timestamps on discover and finish Overlays a forest on the graph Subtree at a node = nodes visited between this node's discovery and finish



DFS algorithm

- → def DFS(G):
 - initialise all vertices to be white, with NULL parent
 - time = 0
 - for u in vertices:
 - if u is white: DFS-Visit(G, u)

→ def DFS-Visit(G,u):

- time++
- u.discovered = time
- u.colour = gray
- for v in u's neighbours:
 - if v is white:
 - v.parent = u
 - DFS-Visit(G, v)
- u.colour = black
- time++
- u.finished = time



why not just call DFS-Visit once?

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Parenthesis structure

- Subtree at a node is visited between the node's discovery and finish times
- Print a "(" when we discover a node u, and ")" when we finish it:
 - Output will be a valid parenthesisation
 - e.g., $(_{u} (_{v} (_{w})_{w})_{v} (_{x} (_{y})_{y})_{x})_{u} (_{z})_{z}$
 - but not: $(_{u} (_{v})_{u})_{v}$
- The (discover, finish) intervals for two vertices are either:
 - Completely disjoint, or
- One contained in the other



Edge classification

Edges in a graph are either:

- Tree edges: in the DFS forest
- Back edges: from a node to an ancestor in the same DFS tree (including self-loop)
- Forward edges: from a node to a descendant

 Cross edges: between nodes in different subtrees or different DFS trees

> Lemma (22.11): For directed graphs, acyclic \iff no back edges

Topological sort

 Linear ordering of vertices such that if u → v is an edge, then u comes before v in sort
Assumes no cycles! (DAG: directed acyclic)
Applications: dependency resolution, compiling files, task planning / Gannt chart

Tweak DFS: as each vertex is finished, insert it at the head of a linked list

e.g.: z, u, x, y, v, w

DFS might not be unique, so topo sort might not be unique



// 1/10

6/9

7/8

Strongly-connected component

Largest completely-connected set of vertices:

 Every vertex in the component has a path to every other vertex in the component

Algorithm:

- Compute DFS(G) to find finishing times
- Let G^T (transpose) be G with all edges reversed
- Compute DFS(G^T) but iterate over neighbours in decreasing order of finishing time from step 1

Each tree in DFS(G^T) is a separate component

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Semester overview



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ch1-4: Intro/definitions (complexity, recurrences, divide-conquer) ch6-8,11: Sorting comparison sorts (insertion, merge, heap, quick) linear sorts (counting, radix, bucket), hash tables ch10,12,18: Data Structures (linked lists, stacks/queues, trees, BST, B-trees) ch15,16: Algorithms (dynamic programming, greedy) ch22: Graph algorithms (BFS, DFS, topo sort, components)



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Exam1: ch1-4

■ Algorith. complexity: $\Theta(=)$, $\Omega(\leq)$, $\Omega(\geq)$, o(<), $\omega(>)$

- Know their technical definitions!
- Proofs!

Solving recurrences: induction, master method
Algorithms to be familiar with:

- Insertion sort, bubble, merge, max subarray
- Matrix multiply (3 algorithms!)



Exam2: ch6, 7

Hand-simulation, complexity analysis
"What if?" questions: tweaks to std algorithms
Ch6: Heapsort

• Trees

• Max heaps: max-heap property, heapify()

- Heapsort: building a heap, using it for sorting
- Priority queue: ops, complexity

Ch7: Quicksort

- Naive quicksort with fixed pivot
- Randomised pivot

Complexity analysis: expected running time E[]

Exam2: ch8, 11

Ch8: Linear-time sorts (assumptions!) • Decision tree model, why $\Omega(n \mid g \mid n)$ comparisons • Counting sort (census + move): $\Theta(n + k)$ • Radix sort (with r-bit digits): $\Theta(d(n + k))$ • Bucket sort: $\Theta(n)$ expected time Ch11: Hash tables Hash function, hash collisions, chaining • Load factor $\alpha = n/(\# \text{ buckets})$, search in $\Theta(1+\alpha)$ Hashes: div, mul, universal hashing Open addressing: linear, guad, double-hash



Exam3: ch10, 12, 18, 15

ch10: Linked-lists (dbl, circ), stacks/queues Implementation and complexity ch12: Trees (terms, expression trees) BSTs (traversal, search, insert, del) ch18: B-trees (motivation, design, variants B*, B+) Operations: search, insert, del • Complexity analysis: CPU and disk ch15: Dynamic programming • Optimal substructure \Rightarrow bottom-up solution Rod-cutting, Fib, matrix-chain, optimal wt BST



Last chapters: ch16, 22

ch16: Greedy algorithms Activity selection Fractional knapsack problem • Huffman coding ch22: Graph algorithms Breadth-first search Depth-first search Edge types, finding cycles • Topological sort • Finding strongly-connected components

